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# ON THE LOCATING CHROMATIC NUMBER OF SUBDIVISION OF BARBELL GRAPHS CONTAINING GENERALIZED PETERSEN GRAPH 

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#### Abstract

The locating chromatic number of a graph is the minimal color required so that it qualifies for some locating coloring. This paper will discuss about the locating chromatic number for the subdivision of barbell graph containing Petersen Graph.


Keyword: locating chromatic number, barbell graph, subdivision, Petersen graph.

## 1. Introduction

The locating chromatic number of a graph was firstly studied by Chartrand et al. [1] as some development of the concept of partition dimension[2] and graph coloring. Consider $G=(V, E)$ as the given connected graph and $c$ as the proper coloring of $G$ using k colors $1,2, \ldots, k$ for some positive integer $k$. We denote $\Pi=\left\{C_{1}, C_{2}, \ldots, C_{k}\right\}$ as the partition of $V(G)$, where $C_{i}$ is the color class, i.e the set of vertices given the $i$-th color, for $i \in[1, k]$. For an arbitrary vertex $v \in V(G)$, the color $\operatorname{code} c_{\Pi}(v)$ is defined as the ordered $k$-tuple

$$
c_{\pi}(v)=\left(d\left(v, C_{1}\right), d\left(v, C_{2}\right), \ldots, d\left(v, C_{k}\right)\right)
$$

where $d\left(v, C_{i}\right)=\min \left\{d(v, x) \mid x \in C_{i}\right\}$ for $i \in[1, k]$. If for every two vertices $u, v \in V(G)$, their color codes are different, $c_{\pi}(u) \neq c_{\pi}(v)$, then $c$ is defined as the locating coloring of $G$ using $k$ colors. The locating chromatic number of $G$, denoted by $\chi_{L}(G)$, is the minimum $k$ such that $G$ has some locating coloring.

There were some interesting results related to the determination of the locating chromatic number of some graphs. The results were obtained by focusing on some certain classes of graph. Chartrand et al. [3] has succeeded in constructing tree on $n$ vertices, $n \geq 5$ with locating chromatic numbers varying from 3 to $n$, except for ( $n-1$ ). Moreover, Asmiati et al. [4] determined the locating chromatic number of homogeneous amalgamation of stars and their monotonicity properties. Recently, Behtoei and Omoomi [5] have obtained the locating chromatic number of the Kneser graph. Asmiati et al. [6] determined the locating chromatic number of Petersen graph and Syofyan et al. [7] trees with certain locating chromatic number.

The barbell graph is constructed by connecting two arbitrary connected graphs $G$ and $H$ by a bridge. Let $B_{P_{n, 1}}$ for $n \geq 3$, be the barbell graph where $G$ and $H$ are two
copies of generalized Petersen graphs $P_{n, 1}$. The following definition of generalized Petersen graph is taken from [8]. Let $\left\{u_{1}, u_{2}, \ldots, u_{n}\right\}$ be the set of vertices in the outer cycle and $\left\{v_{1}, v_{2}, \ldots, v_{n}\right\}$ be the set of vertices in the inner cycle of the Petersen graph, for $\mathrm{n} \geq 3$. Denote the generalized Petersen graph by $P_{n, k}$. From the definition, it is clear that for $n \geq 3$ and $1 \leq k \leq[n-1) / 2]$, the generalized Petersen graph has $2 n$ vertices and $3 n$ edges.

In [9], the locating chromatic number of the barbell graph containing two copies of generalized Petersen graphs $P_{n, 1}$ has been obtained in the following theorem.

Theorem 1.1 [9] For $n \geq 3$, the locating chromatic number of barbell graph $B_{P_{n, 1}}$ is 4 for odd $n$ and 5 otherwise.

This paper will determine the locating chromatic of some graph constructed by subdividing the bridge of the barbell that contains the generalized Petersen graph, denoted by $B_{P_{n, 1}}^{* s}$. This problem is inspired by the results of research Purwasih et. al [10] about the locating chromatic number for a subdivision of a graph on one edge.

## 2. Results and Discussion

In the following theorem, it is discussed about the locating chromatic number for subdivision of some barbell graph containing Petersen graph, denoted by $B_{P_{n, 1}}^{* s}$.

## Theorem 2.1

Let $B_{P_{n, 1}}^{* s}$ be a subdivision of barbell graph containing Petersen Graph for $s \geq 1$. Then the locating chromatic number of $B_{P_{n, 1}}^{* s}$ is 4 for odd $n, n \geq 3$ or 5 for $n$ even, $n \geq 4$.

Proof. Let $B_{P_{n, 1}}^{* s}$ be a subdivision barbell graph for $n \geq 3, s \geq 1$ with the vertex set
$V\left(B_{P_{n, 1}}^{* s}\right)=\left\{u_{i}, u_{n+i}, w_{i}, w_{n+i}: 1 \leq i \leq n\right\} \cup\left\{v_{i}: 1 \leq i \leq s\right\}$, and edge set
$E\left(B_{P_{n, 1}}^{* s}\right)=\left\{u_{i} u_{i+1}, u_{n+i} u_{n+i+1}, w_{i} w_{i+1}, w_{n+i} w_{n+i+1}: 1 \leq i \leq n-1\right\} \cup$
$\left\{u_{n} u_{1}, u_{2 n} u_{n+1}, w_{n} w_{1}, w_{2 n} w_{n+1}\right\} \cup\left\{u_{i} u_{n+1}, w_{i} w_{n+1}: 1 \leq i \leq n\right\} \cup$
$\left\{u_{n} w_{n}\right\} \cup\left\{u_{n} v_{1}, v_{s} w_{n}\right\} \cup\left\{v_{i} v_{i+1}: 1 \leq i \leq s-1\right\}$.

Let us distinguish four cases.
Case 1. $n$ odd. According to Theorem 1.1, it is clear that $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \geq 4$.
To determine the upper bound for the locating chromatic number of subdivision Petersen graph $B_{P_{n, 1}}^{* s}$, construct some locating coloring $c$ using 4 colors as follows.

For odd $s$, define the following coloring

$$
c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for even } i, i \geq 2 \\ 4 & , \text { for odd } i, i \geq 3\end{cases}
$$

$$
c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\ 3 & , \text { for odd } i, i \geq 3 \\ 4 & , \text { for even } i, i \geq 2\end{cases}
$$

$$
c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-2 \\ 2 & , \text { for even } i, i \leq n-1 \\ 4 & , \text { for } i=n\end{cases}
$$

$$
c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-1 \\ 2 & , \text { for odd } i, i \leq n-2 \\ 3 & , \text { for } i=n\end{cases}
$$

$$
c\left(v_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for odd } i, i \geq 3 \\ 4 & , \text { for even } i, i \geq 2\end{cases}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for odd $n$ and $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 1^{\text {st }} \text { component, } i>\frac{n+1}{2} \\ n-i+2 & , \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 3^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 4^{\text {th }} \text { component, } i \text { iodd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ i-1 & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 4^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(w_{i}\right)= \begin{cases}i & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\ i+1 & , \text { for } 3^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\ n-i & , \text { for } 4^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\ n-i+1, \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2} \\ 0 & , \text { for } 2^{n d} \text { component, } i \text { even, } i \leq n-1 \\ & \text { for } 1^{s t} \text { component, } i \text { odd, } i \leq n-2 \\ 1 & , \text { otherwise } .\end{cases}
$$

$$
\begin{aligned}
& c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 3^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\
i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\
n-i & , \text { for } 3^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\
n-i+1, \text { for } 4^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\
0 & , \text { for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-1 \\
& \text { for } 2^{n d} \text { component, } i \text { odd, } i \leq n-2 \\
1 & , \text { otherwise. }\end{cases} \\
& c_{\Pi}\left(v_{i}\right)= \begin{cases}i-1 & , \text { for } 1^{s t} \text { component, } i \leq \frac{s+1}{2} \\
i+2 & , \text { for } 2^{\text {nd }} \text { component, } i \leq \frac{s-1}{2} \\
s-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{s+1}{2} \\
& \text { for } 2^{\text {nd }} \text { component, } i \geq \frac{s+1}{2} \\
2 & , \text { for } 3^{\text {th }} \text { component, } i=1 \\
0 & , \text { for } 4^{t h} \text { component, } i \text { even, } i \geq 2 \\
& \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3 \\
1 & , \text { otherwise. }\end{cases}
\end{aligned}
$$

For even $s$, define the following coloring.

$$
\begin{aligned}
& c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, i \geq 2 \\
4 & , \text { for odd } i, i \geq 3\end{cases} \\
& c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\
3 & , \text { for odd } i, i \geq 3 \\
4 & , \text { for even } i, i \geq 2\end{cases}
\end{aligned}
$$

$$
\begin{aligned}
& c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-2 \\
2 & , \text { for even } i, i \leq n-1 \\
4 & , \text { for } i=n\end{cases} \\
& c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-1 \\
2 & , \text { for odd } i, i \leq n-2 \\
3 & , \text { for } i=n .\end{cases} \\
& c\left(v_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, i \geq 2 \\
4 & , \text { for odd } i, i \geq 3 .\end{cases}
\end{aligned}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for odd $n$ and even $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ i-1 \quad, \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 1^{\text {st }} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 3^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 4^{\text {th }} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{s t} \text { component, } i \leq \frac{n+1}{2} \\ i-1 \quad & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 1^{\text {st }} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 4^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3\end{cases}
$$

$$
\begin{aligned}
& 1 \text {, otherwise. } \\
& c_{\Pi}\left(w_{i}\right)= \begin{cases}i & , \text { for } 4^{t h} \text { component, } i \leq \frac{n-1}{2} \\
i+1 & , \text { for } 3^{t h} \text { component, } i \leq \frac{n-1}{2} \\
n-i & , \text { for } 4^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\
n-i+1, \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2} \\
0 & , \text { for } 2^{n d} \text { component, } i \text { even, } i \leq n-1 \\
1 & , \text { for } 1^{s t} \text { component, } i \text { odd, } i \leq n-2 \\
1 & \end{cases} \\
& c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 3^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\
i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\
n-i & , \text { for } 3^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\
n-i+1, \text { for } 4^{\text {th }} \text { component, } i \geq \frac{n+1}{2} \\
0 & \text {,for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-1 \\
& \text { for } 2^{\text {nd }} \text { component, } i \text { odd, } i \leq n-2 \\
1 & , \text { otherwise. }\end{cases} \\
& c_{\Pi}\left(v_{i}\right)= \begin{cases}i-1 & , \text { for } 1^{s t} \text { component, } i \leq\left(\frac{s}{2}\right)+1 \\
i+2 & , \text { for } 2^{\text {nd }} \text { component, } i \leq \frac{s}{2} \\
s-i+2 & , \text { for } 1^{s t} \text { component, } i>\left(\frac{s}{2}\right)+1 \\
& \text { for } 2^{\text {nd }} \text { component, } i>\frac{s}{2} \\
2 & , \text { for } 3^{\text {th }} \text { component, } i=1 \\
0 & , \text { for } 3^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\
& \text { for } 4^{\text {th }} \text { component, } i \text { odd, } i \geq 3 \\
1 & , \text { otherwise }\end{cases}
\end{aligned}
$$

Since all vertices in $B_{P_{n, 1}}^{* s}$ for odd $n$ have distinct color codes, then c is the locating coloring using 4 colors. So, $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \leq 4$.

Case 2 ( $\boldsymbol{n}$ even). By Theorem 1.1, it is clear that have $\chi_{L}\left(B_{P_{n, 1}}\right) \geq 5$. Consider the following two sub cases.
For $s$ odd, let $c$ be a coloring using 5 colors as follows.

$$
\begin{aligned}
& c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, 2 \leq i \leq n-2 \\
4 & , \text { for odd } i, 3 \leq i \leq n-1 \\
5 & , \text { for } i=n\end{cases} \\
& c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\
3 & , \text { for odd } i, i \geq 3 \\
4 & , \text { for even } i, i \geq 2\end{cases} \\
& c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-1 \\
2 & , \text { for even } i, i \leq n-2 \\
3 & , \text { for } i=n\end{cases}
\end{aligned}
$$

$$
c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-2 \\ 2 & , \text { for odd } i, i \leq n-1 \\ 4 & , \text { for } i=n\end{cases}
$$

$$
c\left(v_{i}\right)= \begin{cases}4 & , \text { for odd } i, 1 \leq i \leq s \\ 5 & , \text { for even } i, 1 \leq i \leq s-1\end{cases}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for even $n$ and odd $s$ are

$$
\begin{aligned}
& c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{\text {nd }} \text { componentand5 }{ }^{t h}, i \leq \frac{n}{2} \\
i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n}{2} \\
n-i & , \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\
n-i+1, & \text { for } 1^{\text {st component, } i>\frac{n}{2}} \\
n-i+2, & \text { for } 2^{n d} \text { component, } i>\frac{n}{2} \\
0 & , \text { for } 3^{t h} \text { component, } i \text { even, } 2 \leq i \leq n-2 \\
2 & \text { for } 4^{t h} \text { component, } i \text { odd, } 2 \leq i \leq n-1 \\
1 & , \text { for } 4^{t h} \text { component, } i=1 \\
\text { for } 3^{t h} \text { component, } i=n \\
& , \text { otherwise. }\end{cases} \\
& 2 \\
& c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n}{2} \\
i-1 & , \text { for } 2^{\text {nd }} \text { component, } i \leq \frac{n}{2} \\
i+1 & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\
n-i+1, & \text { for } 2^{n d} \text { component and } 5^{t h}, i>\frac{n}{2} \\
n-i+2 & , \text { for } 1^{s t} \text { component, } i>\frac{n}{2} \\
0 & , \text { for } 3^{t h} \text { component, } i \text { odd, } 3 \leq i \leq n-1 \\
2 & \text { for } 4^{\text {th }} \text { component, } i \text { even, } 2 \leq i \leq n \\
1 & , \text { for } 3^{t h} \text { component, } i=1\end{cases}
\end{aligned}
$$

$$
c_{\Pi}\left(w_{i}\right)= \begin{cases}i & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n}{2} \\ & \text { for } 3^{\text {th }} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\ n-i & , \text { for } 5^{\text {th }} \text { component, } i>\frac{n}{2} \\ n-i+1, & \text { for } 4^{\text {th }} \text { component, } i>\frac{n}{2} \\ n-i-1, & \text { for } 3^{\text {th }} \text { component, } \frac{n}{2} \leq i \leq n-1 \\ 0 & , \text { for } 1^{s t} \text { component, } i \text { odd, } i \leq n-1 \\ & \text { for } 2^{n d} \text { component, } i \text { even, } i \leq n-2 \\ 2 & , \text { for } 1^{s t} \text { component, } i=n-1 \\ 1 & \text { for } 2^{n d} \text { component, } i=n \\ & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 5^{\text {th }} \text { component, } i \leq \frac{n}{2} \\ i+2 & , \text { for } 3^{\text {th }} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\ n-i & , \text { for } 3^{\text {th }} \text { component, } \frac{n}{2} \leq i \leq n-1 \\ & \text { for } 4^{\text {th }} \text { component, } i>\frac{n}{2} \\ n-i+1, \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-2 \\ & \text { for } 2^{n d} \text { component, } i \text { odd, } i \leq n-1 \\ 2 & , \text { for } 1^{\text {st }} \text { component and } 3^{t h}, i=n \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(v_{i}\right)= \begin{cases}i+1 & , \text { for } 1^{s t} \text { component, } i \leq \frac{s+1}{2} \\ i+2 & , \text { for } 2^{\text {nd }} \text { component, } i \leq \frac{s+1}{2} \\ & \text { for } 3^{t h} \text { component, } i \leq \frac{s-1}{2} \\ s-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{s+1}{2} \\ & \text { for } 3^{t h} \text { component, } i \geq \frac{s+1}{2} \\ s-i+3, & \text { for } 2^{\text {nd }} \text { component, } i>\frac{s+1}{2} \\ 0 \quad & \text { for } 4^{t h} \text { component, } i \text { odd, } 1 \leq i \leq s \\ & \text { for } 5^{t h} \text { component, } i \text { even, } 1 \leq i \leq s-1 \\ 1 & , \text { otherwise. }\end{cases}
$$

For sub case $s$ even, we have

$$
\begin{aligned}
& c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, 2 \leq i \leq n-2 \\
4 & , \text { for odd } i, 2 \leq i \leq n-1 \\
5 & , \text { for } i=n\end{cases} \\
& c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\
3 & , \text { for odd } i, i \geq 3 \\
4 & , \text { for even } i, i \geq 2\end{cases} \\
& c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-1 \\
2 & , \text { for } i=n-1\end{cases} \\
& 3
\end{aligned}
$$

$$
\begin{aligned}
& c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-2 \\
2 & , \text { for odd } i, i \leq n-1 \\
5 & , \text { for } i=n .\end{cases} \\
& c\left(v_{i}\right)= \begin{cases}4 & , \text { for odd } i, 1 \leq i \leq s-1 \\
5 & , \text { for even } i, 1 \leq i \leq s\end{cases}
\end{aligned}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for even $n$ and $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { componentand } 5^{t h}, i \leq \frac{n}{2} \\ i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n}{2} \\ n-i & , \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\ n-i+1, \text { for } 1^{\text {st }} \text { component, } i>\frac{n}{2} \\ n-i+2, \text { for } 2^{n d} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 3^{\text {th }} \text { component, } i \text { even, } 2 \leq i \leq n-2 \\ & \text { for } 4^{\text {th }} \text { component, } i \text { iodd, } 2 \leq i \leq n-1 \\ 2 & , \text { for } 4^{t h} \text { component, } i=1 \\ \text { for } 3^{t h} \text { component, } i=n \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{s t} \text { component, } i \leq \frac{n}{2} \\ i-1 & , \text { for } 2^{n d} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ n-i+1, \text { for } 2^{n d} \text { componentand } 5^{t h}, i>\frac{n}{2} \\ n-i+2, \text { for } 1^{\text {st }} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 3^{t h} \text { component, } i \text { odd, } 2 \leq i \leq n-1 \\ & \text { for } 4^{t h} \text { component, } i \text { even, } 2 \leq i \leq n \\ 2 & , \text { for } 3^{t h} \text { component, } i=1 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & \text {,for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n}{2} \\ i+2 & \text {,for } 3^{\text {th }} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\ n-i & , \text { for } 3^{\text {th }} \text { component, } \frac{n}{2} \leq i \leq n-1 \\ & \text { for } 5^{\text {th }} \text { component, } i>\frac{n}{2} \\ n-i+1, \text { for } 4^{\text {th }} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-2 \\ \text { for } 2^{\text {nd }} \text { component, } i \text { odd, } i \leq n-1 \\ 2 & , \text { for } 1^{\text {st component and } 3^{t h}, i=n} \\ 1 & , \text { otherwise. }\end{cases}
$$

Since all vertices in $B_{P_{n, 1}}^{* s}$ for even $n$ have distinct color codes, then c is the locating coloring using 4 colors. Therefore, the locating chromatic number of the subdividing barbell graph containing generalized Petersen graph, $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \leq$ 5.This completes the proof.

## 2.Acknowledgement

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# On the Locating Chromatic Number of Subdivision of Barbell Graphs Containing Generalized Petersen Graph 

Comments and Suggestions for Authors

In this paper the authors have determined the locating-chrormatic number of Subdivision of Barbell Graphs Containing Generalized Petersen Graphs. The main results is given one big theorem. I have the following comments about the paper. Can the author give some examples of Subdivision of barbell graphs containing generalized graphs for small values of n?. The current version of the paper is very dificult to read and follow. I hope few visual resuts wil make more readable. It would be very helpful if the author can demonstrate theorem with the help of small the graphs (at least for one case). From the title and the introduction of the manuscript, one gets the feeling that the authors have proposed a procedure (algorithm) for determining the locatingchromatic number. However, according to Theorem the locating number of the graphs is 4 or 5 . $\mathrm{d}(\mathrm{x}, \mathrm{y})$ is not defined in the paper.

# On the Locating Chromatic Number of Subdivision of Barbell Graphs Containing Generalized Petersen Graph 

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#### Abstract

The locating chromatic number of a graph is the minimal color required so that it qualifies for some locating coloring. This paper will discuss about the locating chromatic number for the subdivision of barbell graph containing Petersen Graph.


## Key words:

locating chromatic number, barbell graph, subdivision, Petersen graph.

## 1. Introduction

The locating chromatic number of a graph was firstly studied by Chartrand et al. [1] as some development of the concept of partition dimension[2] and graph coloring. Consider $G=(V, E)$ as the given connected graph and $c$ as the proper coloring of $G$ using k colors $1,2, \ldots, k$ for some positive integer $k$. We denote $\Pi=\left\{C_{1}, C_{2}, \ldots, C_{k}\right\}$ as the partition of $V(G)$, where $C_{i}$ is the color class, i.e the set of vertices given the $i$-th color, for $i \in[1, k]$. For an arbitrary vertex $v \in V(G)$, the color $\operatorname{code} c_{\Pi}(v)$ is defined as the ordered $k$-tuple

$$
c_{\pi}(v)=\left(d\left(v, C_{1}\right), d\left(v, C_{2}\right), \ldots, d\left(v, C_{k}\right)\right)
$$

where $d\left(v, C_{i}\right)=\min \left\{d(v, x) \mid x \in C_{i}\right\}$ for $i \in[1, k]$. If for every two vertices $u, v \in V(G)$, their color codes are different, $c_{\pi}(u) \neq c_{\pi}(v)$, then $c$ is defined as the locating coloring of $G$ using $k$ colors. The locating chromatic number of $G$, denoted by $\chi_{L}(G)$, is the minimum $k$ such that $G$ has some locating coloring.
There were some interesting results related to the determination of the locating chromatic number of some graphs. The results were obtained by focusing on some certain classes of graph. Chartrand et al. [3] has succeeded in constructing tree on $n$ vertices, $n \geq 5$ with locating chromatic numbers varying from 3 to $n$, except for $(n-1)$. Moreover, Asmiati et al. [4] determined the locating chromatic number of homogeneous amalgamation of stars and their monotonicity properties. Recently, Behtoei and

Omoomi [5] have obtained the locating chromatic number of the Kneser graph. Asmiati et al. [6] determined the locating chromatic number of Petersen graph and Syofyan et al. [7] trees with certain locating chromatic number.
The barbell graph is constructed by connecting two arbitrary connected graphs $G$ and $H$ by a bridge. Let $B_{P_{n, 1}}$ for $n \geq 3$, be the barbell graph where $G$ and $H$ are two copies of generalized Petersen graphs $P_{n, 1}$. The following definition of generalized Petersen graph is taken from [8]. Let $\left\{u_{1}, u_{2}, \ldots, u_{n}\right\}$ be the set of vertices in the outer cycle and $\left\{v_{1}, v_{2}, \ldots, v_{n}\right\}$ be the set of vertices in the inner cycle of the Petersen graph, for $n \geq 3$. Denote the generalized Petersen graph by $P_{n, k}$. From the definition, it is clear that for $n \geq 3$ and $1 \leq k \leq\lceil n-1) / 2\rceil$, the generalized Petersen graph has $2 n$ vertices and $3 n$ edges.
In [9], the locating chromatic number of the barbell graph containing two copies of generalized Petersen graphs $P_{n, 1}$ has been obtained in the following theorem.
Theorem 1.1 [9] For $n \geq 3$, the locating chromatic number of barbell graph $B_{P_{n, 1}}$ is 4 for odd $n$ and 5 otherwise.
This paper will determine the locating chromatic of some graph constructed by subdividing the bridge of the barbell that contains the generalized Petersen graph, denoted by $B_{P_{n, 1}}^{* s}$. This problem is inspired by the results of research Purwasih et. al [10] about the locating chromatic number for a subdivision of a graph on one edge.

## 2. Results and Discussion

In the following theorem, it is discussed about the locating chromatic number for subdivision of some barbell graph containing Petersen graph, denoted by $B_{P_{n, 1}}^{* s}$.

### 2.1 Theorem

Let $B_{P_{n, 1}}^{* s}$ be a subdivision of barbell graph containing Petersen Graph for $s \geq 1$. Then the locating chromatic number of $B_{P_{n, 1}}^{* s}$ is 4 for odd $n, n \geq 3$ or 5 for $n$ even, $n \geq$ 4.

Proof. Let $B_{P_{n, 1}}^{* s}$ be a subdivision barbell graph for $n \geq$ $3, s \geq 1$ with the vertex set
$V\left(B_{P_{n, 1}}^{* s}\right)=\left\{u_{i}, \quad u_{n+i}, w_{i}, w_{n+i}: 1 \leq i \leq n\right\} \cup\left\{v_{i}: 1 \leq\right.$ $i \leq s\} \quad$, and edge set $\quad E\left(B_{P_{n, 1}^{* s}}^{* s}\right)=\left\{u_{i} u_{i+1}\right.$, $\left.u_{n+i} u_{n+i+1}, w_{i} w_{i+1}, w_{n+i} w_{n+i+1}: 1 \leq i \leq n-1\right\} \cup$
$\left\{u_{n} u_{1}, u_{2 n} u_{n+1}, w_{n} w_{1}, w_{2 n} w_{n+1}\right\} \cup\left\{u_{i} u_{n+1}, w_{i} w_{n+1}:\right.$ $1 \leq i \leq n\} \cup$
$\left\{u_{n} w_{n}\right\} \cup\left\{u_{n} v_{1}, v_{s} w_{n}\right\} \cup\left\{v_{i} v_{i+1}: 1 \leq i \leq s-1\right\}$.
Let us distinguish four cases.
Case 1. $n$ odd. According to Theorem 1.1, it is clear that $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \geq 4$.
To determine the upper bound for the locating chromatic number of subdivision Petersen graph $B_{P_{n, 1}}^{* S}$, construct some locating coloring $c$ using 4 colors as follows.
For odd $s$, define the following coloring
$c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for even } i, i \geq 2 \\ 4 & , \text { for odd } i, i \geq 3\end{cases}$
$c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\ 3 & , \text { for odd } i, i \geq 3 \\ 4 & , \text { for even } i, i \geq 2 .\end{cases}$
$c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-2 \\ 2 & , \text { for even } i, i \leq n-1 \\ 4 & , \text { for } i=n .\end{cases}$
$c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-1 \\ 2 & , \text { for odd } i, i \leq n-2 \\ 3 & , \text { for } i=n .\end{cases}$
$c\left(v_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for odd } i, i \geq 3 \\ 4 & , \text { for even } i, i \geq 2 .\end{cases}$
The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for odd $n$ and $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 1^{\text {st }} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 2^{\text {nd }} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 3^{t h} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 4^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ i-1 \quad, \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 1^{\text {st }} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 4^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ & \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(w_{i}\right)= \begin{cases}i & , \text { for } 4^{t h} \text { component, } i \leq \frac{n-1}{2} \\ i+1 & , \text { for } 3^{t h} \text { component, } i \leq \frac{n-1}{2} \\ n-i & , \text { for } 4^{t h} \text { component, } i \geq \frac{n+1}{2} \\ n-i+1, & \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2} \\ 0 & , \text { for } 2^{n d} \text { component, } i \text { even, } i \leq n-1 \\ & \text { for } 1^{s t} \text { component, } i \text { odd, } i \leq n-2 \\ 1 & , \text { otherwise } .\end{cases}
$$

$$
c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 3^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\ i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n-1}{2} \\ n-i & , \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2} \\ n-i+1, & \text { for } 4^{t h} \text { component, } i \geq \frac{n+1}{2} \\ 0 & , \text { for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-1 \\ & \text { for } 2^{n d} \text { component, } i \text { odd, } i \leq n-2 \\ 1 & , \text { otherwise. }\end{cases}
$$

$v_{\Pi}\left(v_{i}\right)= \begin{cases}i-1 & , \text { for } 1^{s t} \text { component, } i \leq \frac{s+1}{2} \\ i+2 & , \text { for } 2^{n d} \text { component, } i \leq \frac{s-1}{2} \\ s-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{s+1}{2} \\ & \text { for } 2^{n d} \text { component, } i \geq \frac{s+1}{2} \\ 2 & , \text { for } 3^{t h} \text { component, } i=1 \\ 0 & , \text { for } 4^{t h} \text { component, } i \text { even, } i \geq 2 \\ 1 & , \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3\end{cases}$
For even s , define the following coloring.
$c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for even } i, i \geq 2 \\ 4 & , \text { for odd } i, i \geq 3\end{cases}$
$c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\ 3 & , \text { for odd } i, i \geq 3 \\ 4 & , \text { for even } i, i \geq 2 .\end{cases}$
$c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-2 \\ 2 & , \text { for even } i, i \leq n-1 \\ 4 & , \text { for } i=n .\end{cases}$
$c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-1 \\ 2 & , \text { for odd } i, i \leq n-2 \\ 3 & , \text { for } i=n .\end{cases}$
$c\left(v_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\ 3 & , \text { for even } i, i \geq 2 \\ 4 & , \text { for odd } i, i \geq 3 .\end{cases}$
The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for odd $n$ and even $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 1^{\text {st component, } i>\frac{n+1}{2}} \\ n-i+2, & \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ 0 & , \text { for } 3^{t h} \text { component, } i \text { even, } i \geq 2 \\ \text { for } 4^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{s t} \text { component, } i \leq \frac{n+1}{2} \\ i-1 & \text {, for } 2^{n d} \text { component, } i \leq \frac{n+1}{2} \\ n-i+1, & \text { for } 2^{n d} \text { component, } i>\frac{n+1}{2} \\ n-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{n+1}{2} \\ 0 & \text {, for } 4^{t h} \text { component, } i \text { even, } i \geq 2 \\ \text { for } 3^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
\begin{cases}i & , \text { for } 4^{t h} \text { component, } i \leq \frac{n-1}{2} \\ i+1 & , \text { for } 3^{\text {th }} \text { component }, i \leq \frac{n-1}{2}\end{cases}
$$

$$
c_{\Pi}\left(w_{i}\right)=\left\{\begin{array}{l}
n-i \quad, \text { for } 4^{t h} \text { component, } i \geq \frac{n+1}{2} \\
n-i+1, \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2}
\end{array}\right.
$$

$$
0 \quad \text {, for } 2^{n d} \text { component, } i \text { even, } i \leq n-1
$$

$$
\text { for } 1^{\text {st }} \text { component, } i \text { odd, } i \leq n-2
$$

, otherwise.

$$
c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 3^{t h} \text { component, } i \leq \frac{n-1}{2} \\ i+1 & , \text { for } 4^{t h} \text { component, } i \leq \frac{n-1}{2} \\ n-i & , \text { for } 3^{t h} \text { component, } i \geq \frac{n+1}{2} \\ n-i+1, & \text { for } 4^{t h} \text { component, } i \geq \frac{n+1}{2} \\ 0 & \text {,for } 1^{s t} \text { component }, i \text { even, } i \leq n-1 \\ \text { for } 2^{n d} \text { component, } i \text { odd, } i \leq n-2 \\ 1 & , \text { otherwise. }\end{cases}
$$

$c_{\Pi}\left(v_{i}\right)= \begin{cases}i-1 & , \text { for } 1^{s t} \text { component, } i \leq\left(\frac{s}{2}\right)+1 \\ i+2 & , \text { for } 2^{n d} \text { component, } i \leq \frac{s}{2} \\ s-i+2, & \text { for } 1^{s t} \text { component, } i>\left(\frac{s}{2}\right)+1 \\ & \text { for } 2^{\text {nd }} \text { component, } i>\frac{s}{2} \\ 2 & , \text { for } 3^{t h} \text { component, } i=1 \\ 0 & , \text { for } 3^{\text {th }} \text { component, } i \text { even, } i \geq 2 \\ \text { for } 4^{t h} \text { component, } i \text { odd, } i \geq 3 \\ 1 & , \text { otherwise }\end{cases}$
Since all vertices in $B_{P_{n, 1}}^{* s}$ for odd $n$ have distinct color codes, then c is the locating coloring using 4 colors. So, $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \leq 4$.
Case 2 ( $\boldsymbol{n}$ even). By Theorem 1.1, it is clear that have $\chi_{L}\left(B_{P_{n, 1}}\right) \geq 5$. Consider the following two sub cases. For $s$ odd, let $c$ be a coloring using 5 colors as follows.

$$
\begin{aligned}
& c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, 2 \leq i \leq n-2 \\
4 & , \text { for odd } i, 3 \leq i \leq n-1 \\
5 & , \text { for } i=n\end{cases} \\
& c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\
3 & , \text { for odd } i, i \geq 3 \\
4 & , \text { for even } i, i \geq 2\end{cases} \\
& c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-1 \\
2 & , \text { for even } i, i \leq n-2 \\
3 & , \text { for } i=n-1 \\
5 & , \text { for } i=n .\end{cases} \\
& c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-2 \\
2 & , \text { for odd } i, i \leq n-1 \\
4 & , \text { for } i=n .\end{cases} \\
& c\left(v_{i}\right)= \begin{cases}4 & , \text { for odd } i, 1 \leq i \leq s \\
5 & , \text { for even } i, 1 \leq i \leq s-1\end{cases}
\end{aligned}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for even $n$ and odd $s$ are

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n}{2} \\ i-1 & , \text { for } 2^{n d} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ n-i+1, & \text { for } 2^{n d} \text { component and } 5^{\text {th }}, i>\frac{n}{2} \\ n-i+2, & \text { for } 1^{\text {st component, } i>\frac{n}{2}} \\ 0 & , \text { for } 3^{\text {th }} \text { component, } i \text { odd, } 3 \leq i \leq n-1 \\ & \text { for } 4^{\text {th }} \text { component, } i \text { even, } 2 \leq i \leq n \\ 2 & , \text { for } 3^{t h} \text { component, } i=1 \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(w_{i}\right)= \begin{cases}i & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 4^{t h} \text { component, } i \leq \frac{n}{2} \\ & \text { for } 3^{t h} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\ n-i & , \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\ n-i+1, & \text { for } 4^{t h} \text { component, } i>\frac{n}{2} \\ n-i-1, & \text { for } 3^{t h} \text { component, } \frac{n}{2} \leq i \leq n-1 \\ 0 & , \text { for } 1^{s t} \text { component, } i \text { odd, } i \leq n-1 \\ & \text { for } 2^{n d} \text { component, } i \text { even, } i \leq n-2 \\ 2 & , \text { for } 1^{s t} \text { component, } i=n-1 \\ \text { for } 2^{n d} \text { component, } i=n \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
\begin{aligned}
& c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & , \text { for } 4^{t h} \text { component, } i \leq \frac{n}{2} \\
i+1 & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\
i+2 & , \text { for } 3^{t h} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\
n-i & , \text { for } 3^{t h} \text { component, } \frac{n}{2} \leq i \leq n-1 \\
& \text { for } 4^{\text {th }} \text { component, } i>\frac{n}{2} \\
n-i+1, & \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\
0 & , \text { for } 1^{s t} \text { component, } i \text { even, } i \leq n-2 \\
\text { for } 2^{n d} \text { component, } i \text { odd, } i \leq n-1 \\
2 & , \text { for } 1^{s t} \text { component and } 3^{\text {th }}, i=n \\
1 & , \text { otherwise. }\end{cases} \\
& c_{\Pi}\left(v_{i}\right)= \begin{cases}i+1 & , \text { for } 1^{s t} \text { component, } i \leq \frac{s+1}{2} \\
i+2 \quad & , \text { for } 2^{n d} \text { component, } i \leq \frac{s+1}{2} \\
& \text { for } 3^{\text {th }} \text { component, } i \leq \frac{s-1}{2} \\
s-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{s+1}{2} \\
& \text { for } 3^{t h} \text { component, } i \geq \frac{s+1}{2} \\
s-i+3, & \text { for } 2^{n d} \text { component, } i>\frac{s+1}{2} \\
0 \quad & \text { for } 4^{t h} \text { component, } i \text { odd, } 1 \leq i \leq s \\
& \text { for } 5^{\text {th }} \text { component, } i \text { even, } 1 \leq i \leq s-1 \\
1 & , \text { otherwise. }\end{cases}
\end{aligned}
$$

For sub case $s$ even, we have

$$
\begin{aligned}
& c\left(u_{i}\right)= \begin{cases}1 & , \text { for } i=1 \\
3 & , \text { for even } i, 2 \leq i \leq n-2 \\
4 & , \text { for odd } i, 2 \leq i \leq n-1 \\
5 & , \text { for } i=n\end{cases} \\
& c\left(u_{n+i}\right)= \begin{cases}2 & , \text { for } i=1 \\
3 & , \text { for odd } i, i \geq 3 \\
4 & , \text { for even } i, i \geq 2\end{cases} \\
& c\left(w_{i}\right)= \begin{cases}1 & , \text { for odd } i, i \leq n-1 \\
2 & , \text { for even } i, i \leq n-2 \\
3 & , \text { for } i=n-1\end{cases}
\end{aligned}
$$

$$
\begin{aligned}
& c\left(w_{n+i}\right)= \begin{cases}1 & , \text { for even } i, i \leq n-2 \\
2 & , \text { for odd } i, i \leq n-1 \\
5 & , \text { for } i=n .\end{cases} \\
& c\left(v_{i}\right)= \begin{cases}4 & , \text { for odd } i, 1 \leq i \leq s-1 \\
5 & , \text { for even } i, 1 \leq i \leq s\end{cases}
\end{aligned}
$$

The color codes of $V\left(B_{P_{n, 1}}^{* s}\right)$ for even $n$ and $s$ are

$$
c_{\Pi}\left(u_{i}\right)= \begin{cases}i & , \text { for } 2^{n d} \text { componentand } 5^{t h}, i \leq \frac{n}{2} \\ i-1 & , \text { for } 1^{\text {st }} \text { component, } i \leq \frac{n}{2} \\ n-i & , \text { for } 5^{t h} \text { component, } i>\frac{n}{2} \\ n-i+1, \text { for } 1^{\text {st }} \text { component, } i>\frac{n}{2} \\ n-i+2, \text { for } 2^{n d} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 3^{t h} \text { component, } i \text { even, } 2 \leq i \leq n-2 \\ \text { for } 4^{t h} \text { component, } i \text { odd, } 2 \leq i \leq n-1 \\ 2 & , \text { for } 3^{t h} \text { component, } i=1 \\ \text { for } 3^{t h} \text { component, } i=n \\ 1 & , \text { otherwise. }\end{cases}
$$

$$
c_{\Pi}\left(u_{n+i}\right)= \begin{cases}i & , \text { for } 1^{s t} \text { component, } i \leq \frac{n}{2} \\ i-1 & , \text { for } 2^{n d} \text { component, } i \leq \frac{n}{2} \\ i+1 & , \text { for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\ n-i+1, & \text { for } 2^{n d} \text { componentand } 5^{t h}, i>\frac{n}{2} \\ n-i+2, & \text { for } 1^{s t} \text { component, } i>\frac{n}{2} \\ 0 & , \text { for } 3^{t h} \text { component, } i \text { odd, } 2 \leq i \leq n-1 \\ 2 & , \text { for } 4^{t h} \text { component, } i \text { even, } 2 \leq i \leq n \\ 1 & , \text { otherwise } .\end{cases}
$$

$$
\begin{aligned}
& c_{\Pi}\left(w_{n+i}\right)= \begin{cases}i & \text {,for } 5^{t h} \text { component, } i \leq \frac{n}{2} \\
i+1 & , \text { for } 4^{\text {th }} \text { component, } i \leq \frac{n}{2} \\
i+2 & \text {,for } 3^{t h} \text { component, } i \leq\left(\frac{n}{2}\right)-1 \\
n-i & , \text { for } 3^{t h} \text { component, } \frac{n}{2} \leq i \leq n-1 \\
n-i+1, \text { for } 4^{\text {th }} \text { component, } i>\frac{n}{2} \\
0 & , \text { for } 1^{\text {st }} \text { component, } i \text { even, } i \leq n-2 \\
& \text { for } 2^{\text {nd }} \text { component, } i \text { odd, } i \leq n-1 \\
2 & , \text { for } 1^{s t} \text { component and } 3^{t h}, i=n \\
1 & , \text { otherwise. }\end{cases} \\
& c_{\Pi}\left(v_{i}\right)= \begin{cases}i+1 & , \text { for } 1^{s t} \text { component, } i \leq \frac{o}{2} \\
i+2 \quad & , \text { for } 2^{n d} \text { component, } i \leq \frac{s}{2} \\
& \text { for } 3^{t h} \text { component, } i \leq \frac{s}{2} \\
s-i+2 & , \text { for } 1^{s t} \text { component, } i>\frac{s}{2} \\
& \text { for } 3^{t h} \text { component, } i>\frac{s}{2} \\
s-i+3 & , \text { for } 2^{n d} \text { component, } i>\frac{s}{2} \\
0 \quad & , \text { for } 4^{t h} \text { component, } i \text { odd, } 1 \leq i \leq s-1 \\
& \text { for } 5^{t h} \text { component, } i \text { even, } 1 \leq i \leq s \\
1 & , \text { otherwise. }\end{cases}
\end{aligned}
$$

Since all vertices in $B_{P_{n, 1}}^{* s}$ for even $n$ have distinct color codes, then c is the locating coloring using 4 colors. Therefore, the locating chromatic number of the subdividing barbell graph containing generalized Petersen graph, $\chi_{L}\left(B_{P_{n, 1}}^{* s}\right) \leq 5$.This completes the proof.

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